

Karlsruhe Institute of Technology

Reliable Computing I

Lecture 6: Information Redundancy

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National Research Center of the Helmholtz Association

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Today's Lecture


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- Code, codeword, binary code
- Error detecting and correcting codes
- Hamming distance and codes
- Parity prediction
 - Odd/even parity
 - Basic parity approaches

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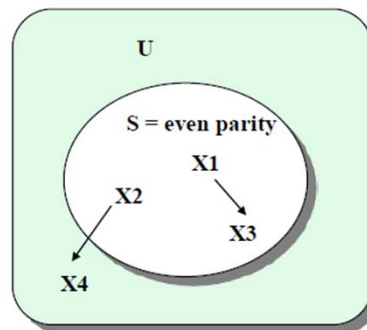
Error Detection through Encoding



- At logic level, codes provide means of masking or detecting errors
- Formally, code is a subset S of universe U of possible vectors
- A noncode word is a vector in set $U-S$

X_1 is a codeword
<10010011>
due to multiple bit error,
becomes
 $X_3 = <10011100>$
not detectable

X_2 is a codeword,
becomes X_4 noncode
detectable



Basic Idea



- Start with k -bit data word
- Add r check bits
- Total = n -bit **codeword** ($n=k+r$)
- Map 2^k data words to 2^n sized codeword space
- Overhead = r/n (sometimes computed as r/k)
 - E.g., for (single-bit) parity, the overhead is $1/n$

Basic Concepts

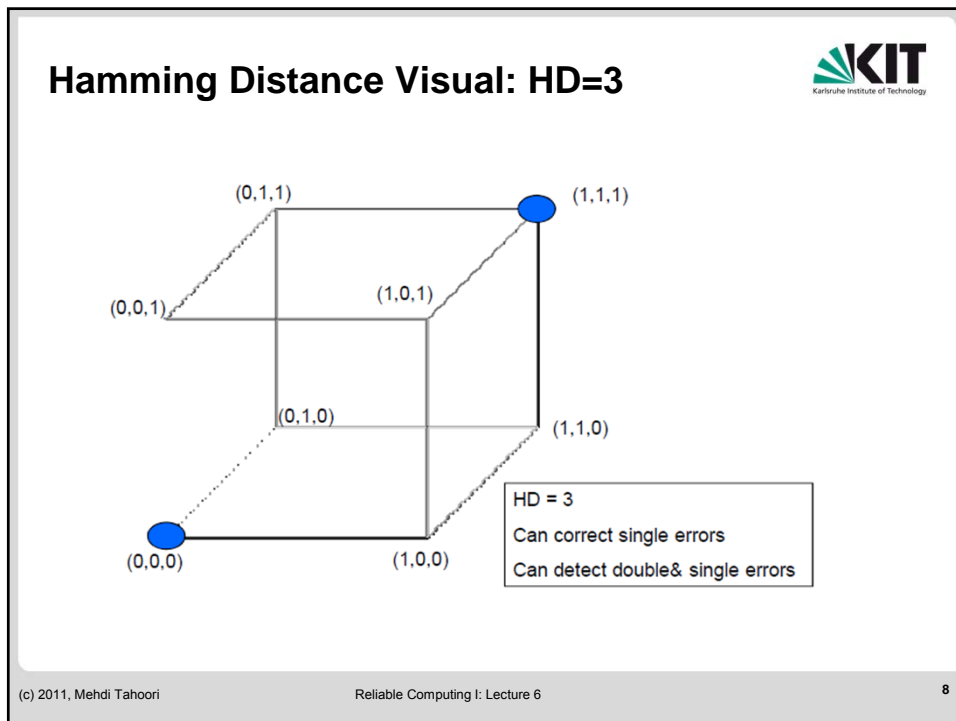
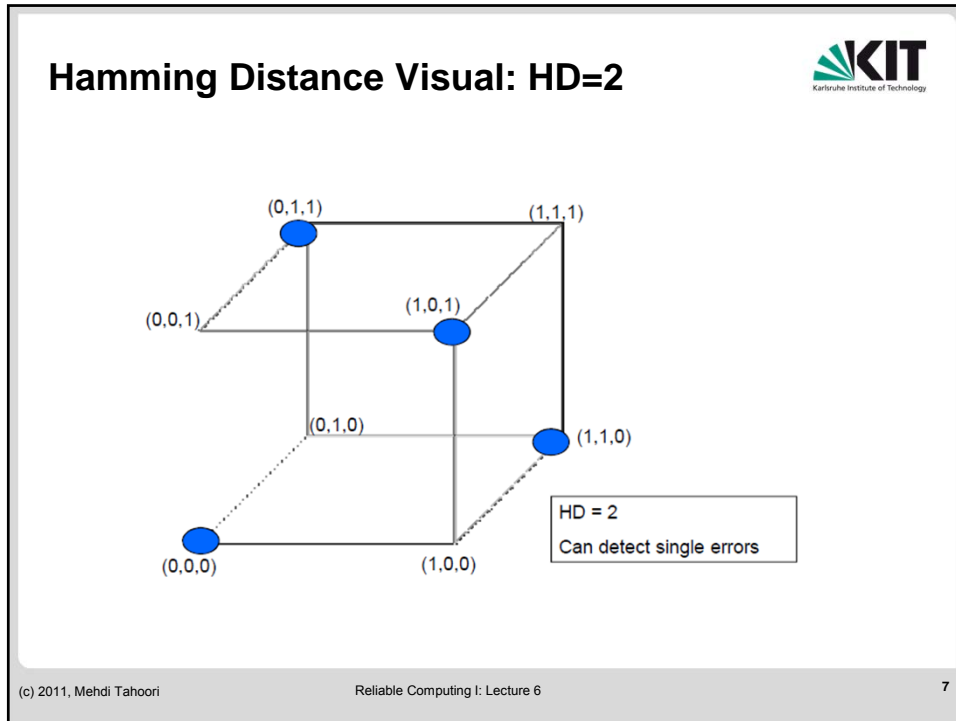


- Code, codeword, encoding, decoding error detection code, error correcting code
- Hamming distance properties:
 - The **Hamming weight** of a vector x (e.g., codeword), $w(x)$, is number of nonzero elements of x .
 - **Hamming distance** between two vectors x and y , $d(x,y)$ is number of bits in which they differ.
 - **Distance of a code** is a minimum of Hamming distances between all pairs of code words.
 - Example: $x = (1011)$, $y = (0110)$
 $w(x) = 3$, $w(y) = 2$, $d(x, y) = 3$

Hamming Distance




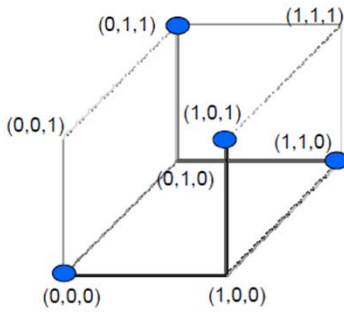
- Hamming distance (HD): number of bits in which two words differ from each other
 - E.g., 0010 and 1110 have a Hamming distance of ??
- For a group of codewords, the minimum HD between any two codewords determines the code's ability to detect and/or correct errors
 - This is a fundamental rule, not just some ad-hoc reasoning



Hamming Distance and Error Detection

■ Can detect up to t -bit errors if $HD \geq t + 1$






HD=2, detects 1-bit errors

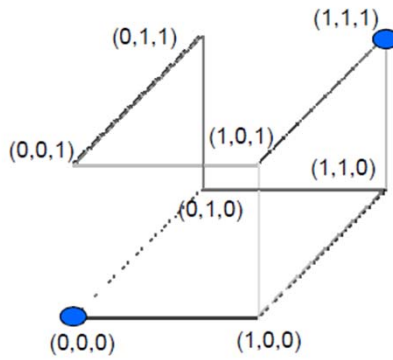
- What if we receive **111**?
 - Could've been **011**
 - Could've been **101**
 - Could've been **110**
 - What about **001**? Or **000**??
- What if we receive **011**?
 - Could it have been **001**???

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Hamming Distance and Error Detection

■ Can detect up to t -bit errors if $HD \geq t + 1$






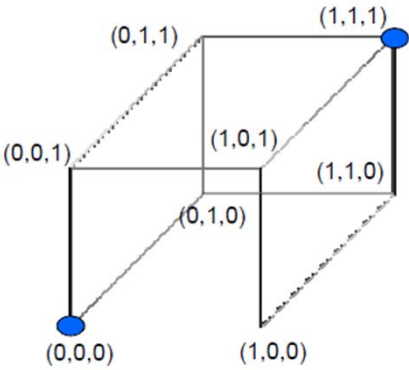
HD=3, detects 1,2-bit errors

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Hamming Distance and Error Correction




■ Can correct up to t -bit errors if $HD \geq 2t+1$

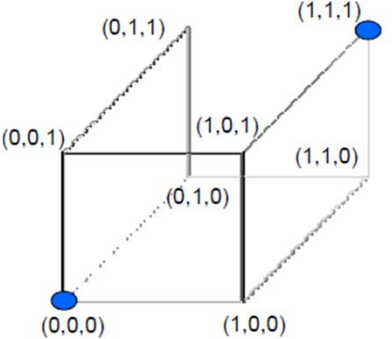


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Hamming Distance and Error Correction



■ Can correct up to t -bit errors if $HD \geq 2t+1$



- **What if we receive 011?**
 - More likely to have been **111**
 - But could've been **000**
 - Guess that it was **111**
- **What if we receive 111?**
 - Could it have been **000**?

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Summary: Hamming Distance Properties




- To **detect** all error patterns of Hamming distance $\leq d$,
code distance must be $\geq d+1$
 - e.g., code with distance 2 can detect patterns with distance 1 (i.e., single bit errors)
- To **correct** all error patterns of Hamming distance $\leq c$,
code distance must be $\geq 2c + 1$
- To correct all patterns of Hamming distance c and detect up to d additional errors ,
code distance must be $\geq 2c + d + 1$
 - e.g., code with distance 3 can detect and correct all single-bit errors

Single-bit Parity



- Simplest error detection code
 - Adds one bit of redundancy to each data word
- Even (odd) parity: add bit such that total number of ones in codeword is even (odd)
 - E.g., 001010 gets a parity bit of 0 for even parity (1 for odd)
- Can detect all single-bit errors
 - Hamming distance ≥ 2
 - Could be greater than 2 if data words don't use all bit combinations
- Drawbacks:
 - Can't detect anything except single-bit errors

Parity Codes - Example



Odd and even parity codes for BCD data

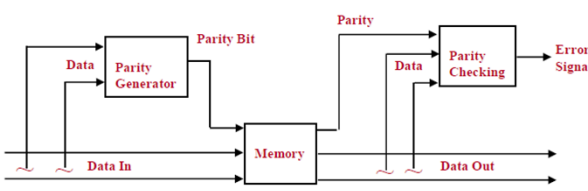
Decimal Digit	BCD	BCD odd parity	BCD even parity
0	0000	00001	00000
1	0001	00010	00011
2	0010	00100	00101
3	0011	00111	00110
4	0100	01000	01001
5	0101	01011	01010
6	0110	01101	01100
7	0111	01110	01111
8	1000	10000	10001
9	1001	10011	10010

↑

Parity Bit

↑


Parity Bit

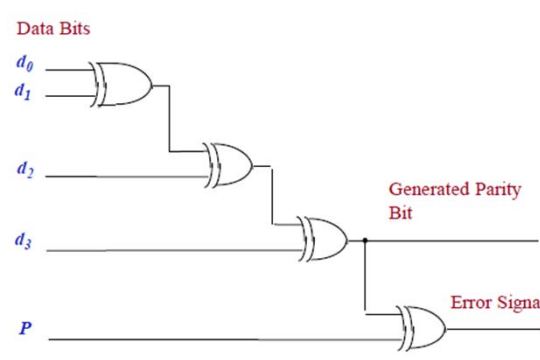


Organization of a memory that uses single-bit parity.
The parity bit is generated when data is written to memory and checked when data is read.

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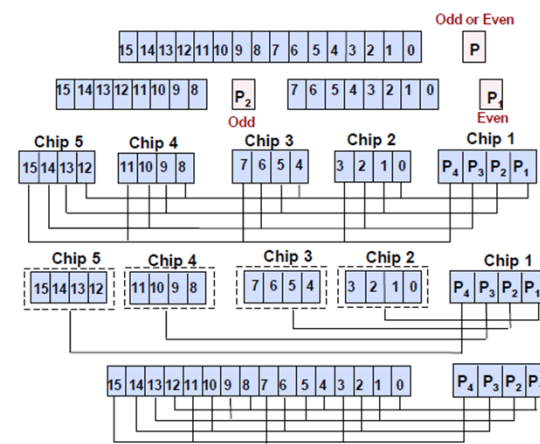
XOR Tree for Parity Generation






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Codes for RAMs



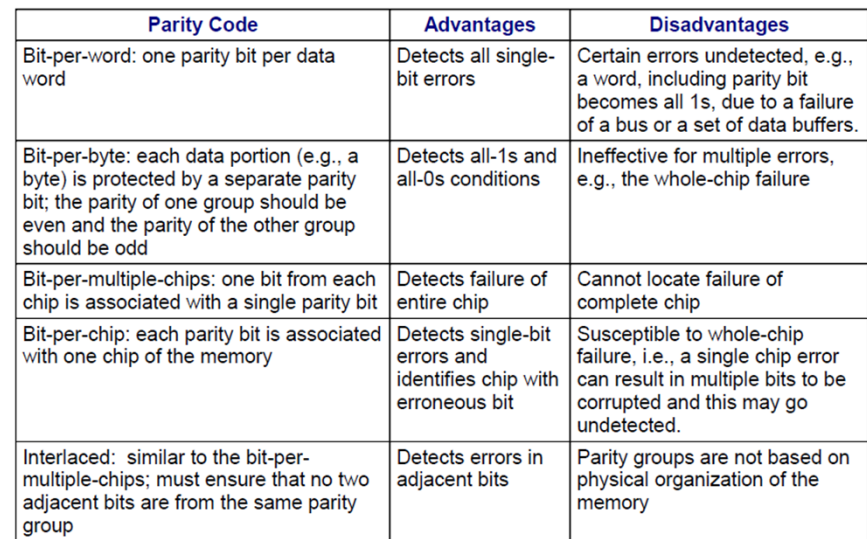
The diagram illustrates five parity coding schemes for RAMs:

- Bit-per-word parity:** A single parity bit P is calculated for the entire word (bits 15 to 0).
- Bit-per-byte parity:** Two parity bits, P_2 (Odd) and P_1 (Even), are calculated for two groups of 8 bits each.
- Bit-per-multiple-chips parity:** A single parity bit P is calculated for all bits across all chips (Chip 5 to Chip 1).
- Bit-per-chip parity:** A parity bit P_i is calculated for each individual chip i .
- Interlaced parity:** A single parity bit P is calculated for bits that are interleaved across all chips.




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Parity Codes for Memory - Comparison



Parity Code	Advantages	Disadvantages
Bit-per-word: one parity bit per data word	Detects all single-bit errors	Certain errors undetected, e.g., a word, including parity bit becomes all 1s, due to a failure of a bus or a set of data buffers.
Bit-per-byte: each data portion (e.g., a byte) is protected by a separate parity bit; the parity of one group should be even and the parity of the other group should be odd	Detects all-1s and all-0s conditions	Ineffective for multiple errors, e.g., the whole-chip failure
Bit-per-multiple-chips: one bit from each chip is associated with a single parity bit	Detects failure of entire chip	Cannot locate failure of complete chip
Bit-per-chip: each parity bit is associated with one chip of the memory	Detects single-bit errors and identifies chip with erroneous bit	Susceptible to whole-chip failure, i.e., a single chip error can result in multiple bits to be corrupted and this may go undetected.
Interlaced: similar to the bit-per-multiple-chips; must ensure that no two adjacent bits are from the same parity group	Detects errors in adjacent bits	Parity groups are not based on physical organization of the memory



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Parity Prediction in Arithmetic Circuits



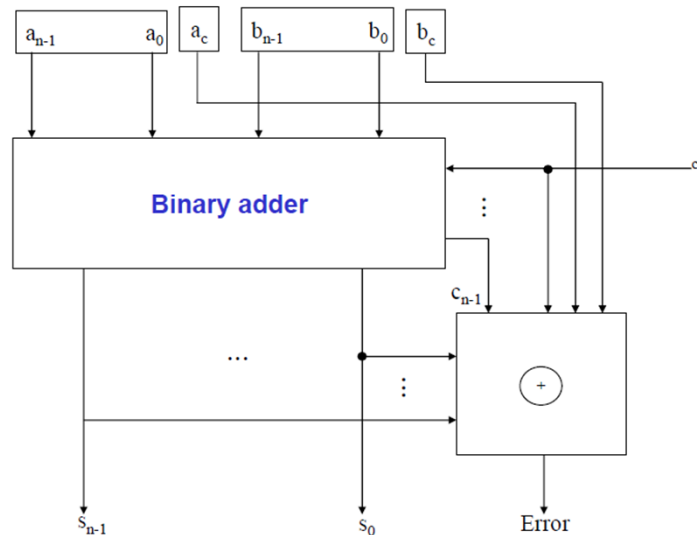
Binary Adder

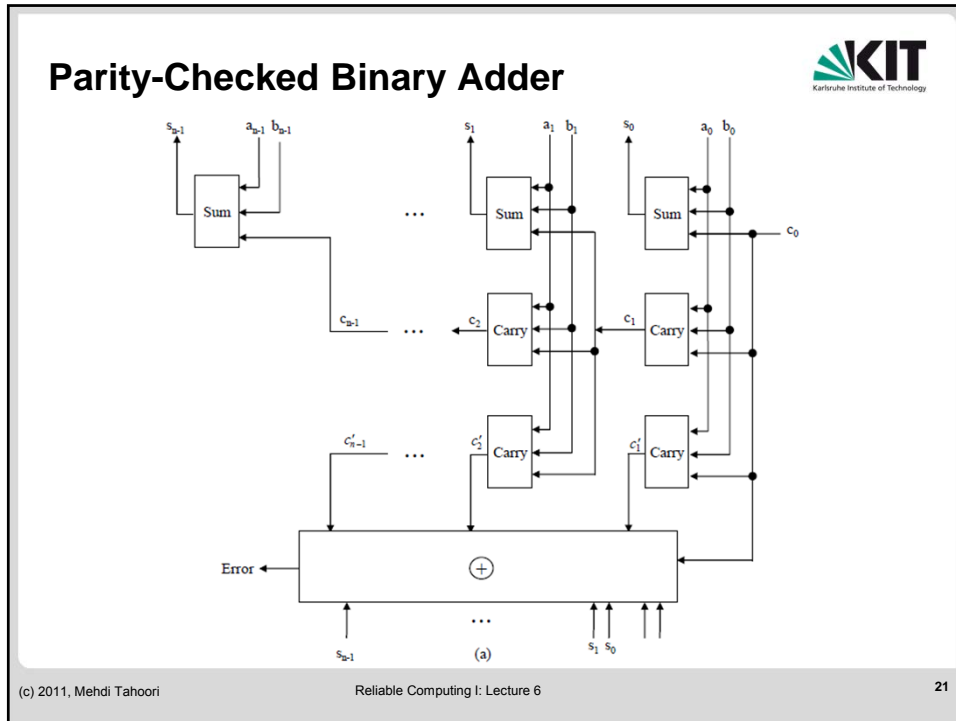
- Two inputs: $a = (a_{n-1} \dots a_0 a_c)$ and $b = (b_{n-1} \dots b_0 b_c)$
- Two operands to be added: $(a_{n-1} \dots a_0)$ and $(b_{n-1} \dots b_0)$
- a_c and b_c are check bits of a and b respectively
- Encoded output will be $s = (s_{n-1} \dots s_0 s_c)$ where $(s_{n-1} \dots s_0)$ are determined by the ordinary binary addition of $(a_{n-1} \dots a_0)$ to $(b_{n-1} \dots b_0)$ and s_c is the check bit for $(s_{n-1} \dots s_0)$

- Then
$$s_c = \sum_{i=0}^{n-1} s_i = \sum_{i=0}^{n-1} a_i \oplus \sum_{i=0}^{n-1} b_i \oplus \sum_{i=0}^{n-1} c_i$$


- Reduces to
$$s_c = a_c \oplus b_c \oplus \sum_{i=0}^{n-1} c_i$$

Parity Prediction in Binary Adder





Binary Multiplier

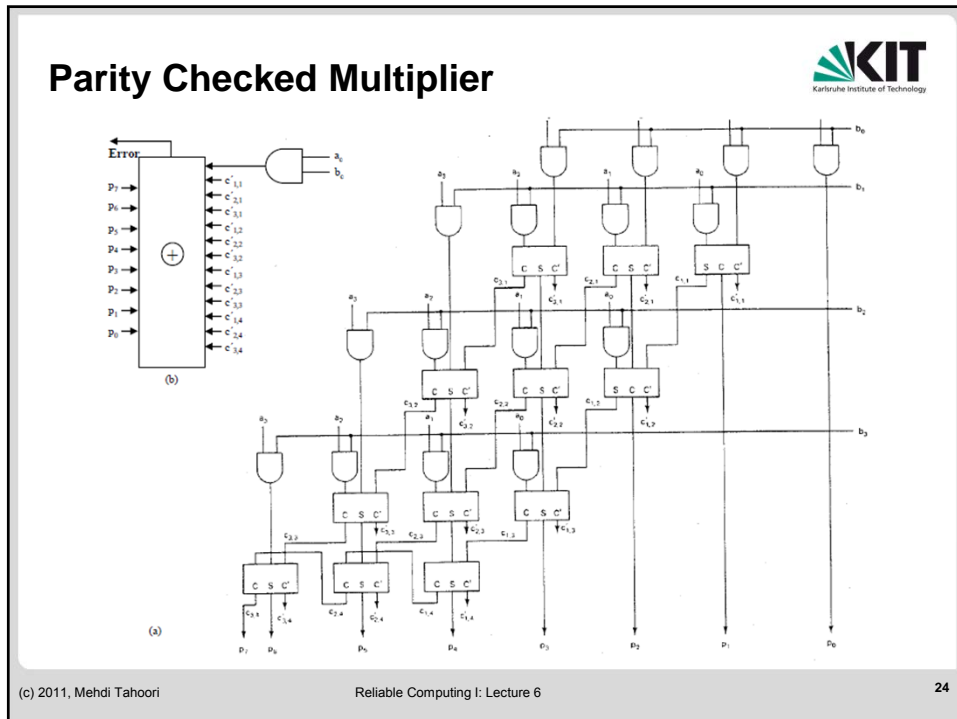
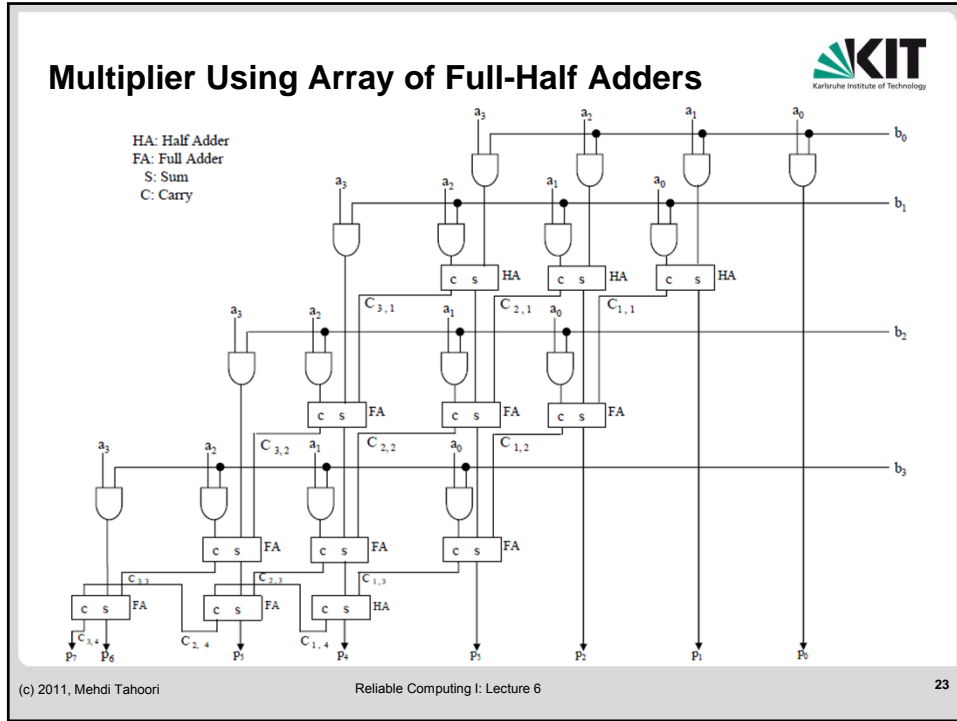


$$\begin{aligned}
 p_0 &= a_0 b_0 \\
 p_1 &= a_0 b_1 \oplus a_1 b_0 \\
 p_2 &= a_0 b_2 \oplus a_1 b_1 \oplus a_2 b_0 \oplus c_{1,1} \\
 p_3 &= a_0 b_3 \oplus a_1 b_2 \oplus a_2 b_1 \oplus a_3 b_0 \oplus c_{2,1} \oplus c_{1,2} \\
 p_4 &= a_1 b_3 \oplus a_2 b_2 \oplus a_3 b_1 \oplus c_{3,1} \oplus c_{2,2} \oplus c_{1,3} \\
 p_5 &= a_2 b_3 \oplus a_3 b_2 \oplus c_{3,2} \oplus c_{2,3} \oplus c_{1,4} \\
 p_6 &= a_3 b_3 \oplus c_{3,3} \oplus c_{2,4} \\
 p_7 &= c_{3,4}
 \end{aligned}$$

■ Therefore, denoting the check bit for $(p_7 \dots p_0)$ by p_c

$$\begin{aligned}
 p_c &= \sum_{i=0}^7 p_i \\
 &= \left(\sum_{i=0}^3 a_i \right) \left(\sum_{i=0}^3 b_i \right) \oplus \sum_{i=1}^3 \sum_{j=1}^4 c_{i,j} \\
 &= a_c b_c \oplus \sum_{i=1}^3 \sum_{j=1}^4 c_{i,j}
 \end{aligned}$$

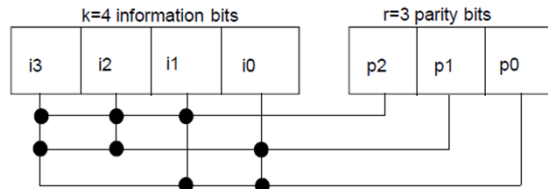
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Overlapping Parity (for single-bit errors)



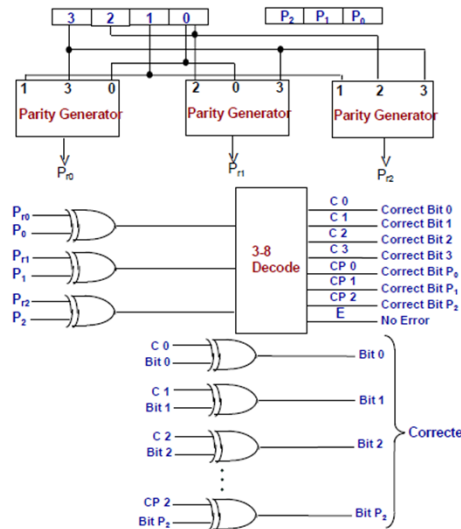
- Parity groups are formed with each bit appearing in more than one parity group
- Errors can be detected and located
- Erroneous bit can be corrected by a simple complementation



Which bit has error?	Parity bits affected
i3	p2, p1, p0
i2	p2, p1
i1	p2, p0
i0	p1, p0
p2	p2
p1	p1
p0	p0

When receiving codeword, re-compute 3 parity bits and compare to those that were sent. If different, can diagnose error (and correct it)!

Error Correction with Overlapped Parity



Generalized Overlapping Parity Codes

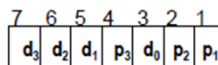


- The previous slide showed how to use overlapping parity to detect and diagnose single-bit errors
- For single-bit errors, there are $k+r$ possible errors
 - Therefore, we need $2^r \geq k + r + 1$ to uniquely diagnose errors
- In general, can extend this scheme to detect and diagnose more than single-bit errors
 - General approach called “Hamming Codes”


Hamming Error-Correcting Code



- Require from 10% to 40% redundancy
- Best thought of as overlapping parity
- The Hamming single-error correcting code uses c parity check bits to protect k bits of information:
 - $2^c \geq c + k + 1$
- Example:
 - suppose four information bits (d_3, d_2, d_1, d_0) and as a result three parity bits (p_1, p_2, p_3)
 - the bits are partitioned into groups as (d_3, d_1, d_0, p_1), (d_3, d_2, d_0, p_2) and (d_3, d_2, d_1, p_3)
 - the grouping of bits can be determine from a list of binary numbers from 0 to $2^k - 1$.
 - each check bit is specified to set the parity, either even or odd, of its respective group



Hamming Error-Correcting Code



Determining the bit groups (three parity bits)

0 0 0			
0 0 1	1		
0 1 0		2	
0 1 1	3	3	
1 0 0			4
1 0 1	5	5	
1 1 0		6	6
1 1 1	7	7	7

Parity bits calculation

p1 = XOR of bits (3, 5, 7)
p2 = XOR of bits (3, 6, 7)
p3 = XOR of bits (5, 6, 7)


Parity checking

c1 = XOR of bits (1, 3, 5, 7)
c2 = XOR of bits (2, 3, 6, 7)
c3 = XOR of bits (4, 5, 6, 7)

■ Observe that each group of bits for parity checking starts with a number that is a power of 2, e.g., 1, 2, 4.

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(7,4) Hamming Code



- Class of (n,k) Hamming codes, e.g., (7,4) [r= n-k =3]
- Let i₁, i₂, i₃, i₄ be the information bits
- Let p₁, p₂, p₄ be the check bits
- p₁ = i₁ xor i₂ xor i₄
- p₂ = i₁ xor i₃ xor i₄
- p₄ = i₂ xor i₃ xor i₄
- Let H be the **Parity Check Matrix**
- If C is a codeword, then H C = 0 (mult modulo 2!)
- Else, H C = S, where S is the **syndrome**
 - Syndrome identifies where error occurred (i.e., which bit)
 - This works out like magic because of some cute math

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(7,4) Hamming Code



■ $\underline{H} =$

p1	p2	i1	p4	i2	i3	i4
1	0	1	0	1	0	1
0	1	1	0	0	1	1
0	0	0	1	1	1	1

■ Info word: 0101: $p_1 = 0$, $p_2 = 1$, $p_4 = 0$

■ codeword is 0100101

■ Example 1:

■ received error-free codeword $R = 0100101$

■ Compute syndrome: $\underline{S} = \underline{H} R = \underline{0} = [0\ 0\ 0]$

■ Example 2:

■ received $R = 0110101$ (i.e., error in bit position 3)

■ Compute syndrome: $S = H R = [1\ 1\ 0]$

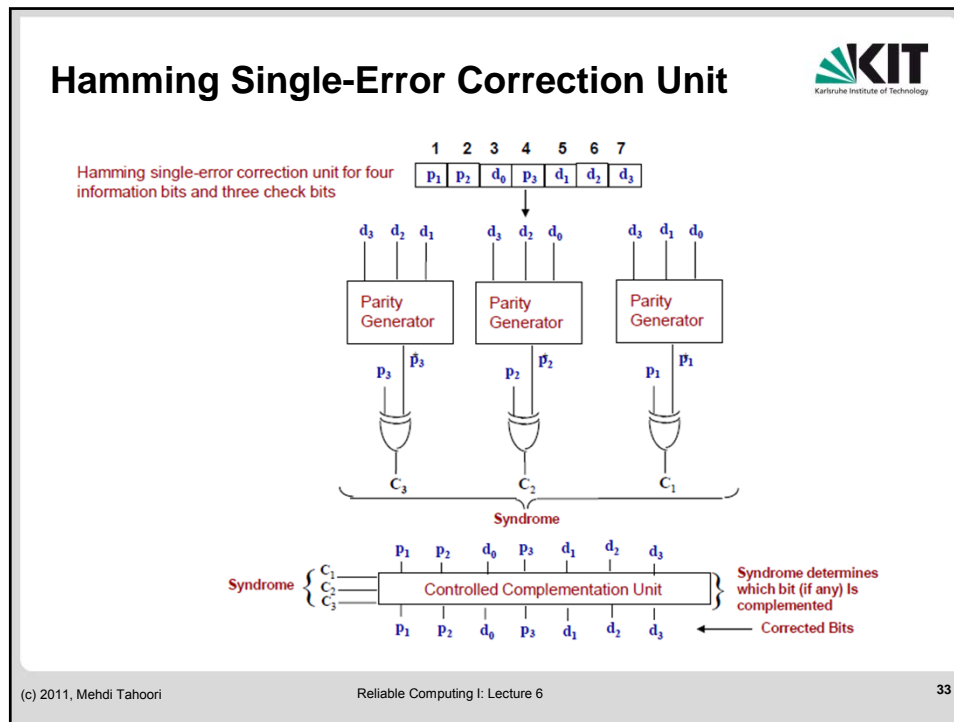
■ Read backwards this is 011 = 3

Check Bits and Syndromes for Single-Bit Errors

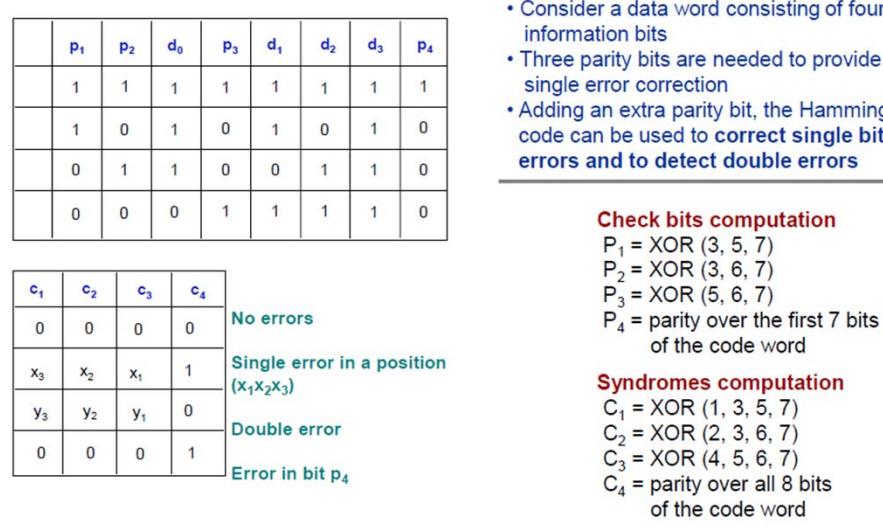


- The original data is encoded by generating a set C_g , of parity bits.
- To check correctness, the encoding process is repeated and a set C_c , of parity bits is generated.
- If C_g and C_c agree, the information is correct.
- If C_g and C_c disagree, the information is incorrect and must be corrected.
- To aid the correction, a *syndrome* is defined:
 - The syndrome is a binary word that has 1 in each bit position in which C_g and C_c disagree; the syndrome points directly to the erroneous bit.

Erroneous bits	Check bits affected	Syndromes
d_0	p_1, p_2	110
d_1	p_1, p_3	101
d_2	p_2, p_3	011
d_3	p_1, p_2, p_3	111
p_1	p_1	100
p_2	p_2	010
p_3	p_3	001



Single Error Correction and Double Error Detection Hamming Code (SEC-DED)



	p_1	p_2	d_0	p_3	d_1	d_2	d_3	p_4
	1	1	1	1	1	1	1	1
	1	0	1	0	1	0	1	0
	0	1	1	0	0	1	1	0
	0	0	0	1	1	1	1	0

c_1	c_2	c_3	c_4	
0	0	0	0	No errors
x_3	x_2	x_1	1	Single error in a position ($x_1x_2x_3$)
y_3	y_2	y_1	0	Double error
0	0	0	1	Error in bit p_4

- Consider a data word consisting of four information bits
- Three parity bits are needed to provide single error correction
- Adding an extra parity bit, the Hamming code can be used to correct single bit errors and to detect double errors

Check bits computation

$P_1 = \text{XOR}(3, 5, 7)$
 $P_2 = \text{XOR}(3, 6, 7)$
 $P_3 = \text{XOR}(5, 6, 7)$
 $P_4 = \text{parity over the first 7 bits of the code word}$

Syndromes computation

$C_1 = \text{XOR}(1, 3, 5, 7)$
 $C_2 = \text{XOR}(2, 3, 6, 7)$
 $C_3 = \text{XOR}(4, 5, 6, 7)$
 $C_4 = \text{parity over all 8 bits of the code word}$

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Single Error Correction and Double Error Detection Hamming Code (SEC-DED) Example



	p_1	p_2	d_0	p_3	d_1	d_2	d_3	p_4
1	1	1	0	0	1	1	0	0
2	1	1	1	0	1	1	0	0
3	1	1	0	0	1	0	0	0
4	0	1	0	0	0	1	0	0
5	1	1	0	0	1	1	0	1

Initial daa
 $d_0 d_1 d_2 d_3$
0 1 1 0

	c_1	c_2	c_3	c_4
1	0	0	0	0
2	1	1	0	1
3	0	1	1	1
4	0	0	1	0
5	0	0	0	1

No errors
Single error in position 3
Single error in position 6
Double error
Error in bit p_4